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ONLINE ALGORITHMS & APPLICATIONS



TOPICS

Online Algorithms
Offline Algorithms
Competitive Analysis
Adversaries
Applications



In online computation a computer algorithm must decide how to act on incoming items of information without any knowledge of future inputs

- How should the next call be routed?
- Which cache block to be removed when the cache is full?



- An <u>online algorithm</u> is one that can process its input piece-by-piece, without having the entire input available from the start
- In contrast, an <u>offline algorithm</u> is given the whole problem data from the beginning and is required to output an answer which solves the problem at hand
- For example, <u>selection sort</u> requires that the entire list be given before it can sort it



- An algorithm is called "online" if it produces (partial) output while still reading its input.
- Some algorithms must be online, because they produce a stream of output for a stream of input; output is produced while the input (which might even be infinite in length) is being read.
- All scheduling algorithms are online algorithms
- When an OS is paging memory, or when a dispatcher is dispatching ambulances around the city, it is often important to be able to guarantee certain levels of performance.
- **OS** or dispatcher have no idea what happens next.
- Must decide strictly according to data available at the time of the action taken



- Input is revealed to the algorithm incrementally
- Output is produced incrementally
- Some output must be produced before the entire input is known to the algorithm
- How to make decisions with partial information?
- Unknown information: the future.



APPLICATIONS

Resource Allocation

- Scheduling
- Memory Management
- Routing

Robot Motion Planning

- Exploring an unknown terrain
- Finding a destination



METHODS OF ANALYSIS

Probabilistic Analysis

- Assume a distribution generating the input.
- Find an algorithm which minimizes the expected cost of the algorithm.
- Pros: can incorporate information predicting the future.
- Cons: can be difficult to determine probability distributions accurately.



METHODS OF ANALYSIS

Competitive Analysis (Worst Case)

- For any input, the cost of our online algorithm is never worse than 'c' times the cost of the optimal offline algorithm.
- Pros: can make very robust statements about the performance of a strategy.
- **Cons: results tend to be pessimistic.**



- Finding a <u>shortest path</u> in a finite connected graph when the <u>graph</u> is unknown and the algorithm receives the <u>node</u> neighbors only when it "enters" the node.
- Problem can not be solved optimally without a simple exhaustive search.
- New performance measures have to be introduced, such as <u>competitive analysis</u>, which compares the performance of an online algorithm with that of a hypothetical offline algorithm that knows the entire input in advance.



- Suppose you decide to learn to ski
- After each trip, you will make an irrevocable decision whether to stop skiing or continue learning
- **You have no idea in advance what your decision will be**
- Skiing is an equipment-intensive sport and before each trip you have two options: rent the equipment at \$x per day or buy the equipment for a grand sum of \$y such that:

y=cx for some integer c>1.

Before each trip to the mountains you have to decide whether to rent or buy

* Example taken from 'An Introduction to Competitive Analysis for Online Optimization' Maurice Queyranne, University of British Columbia.



- OBJECTIVE: to minimize cost
- Buying equipment even before taking one lesson would be a terrible waste if you decide to stop after the first trip
- On the other hand, if you take many trips then at some point it would be cheaper to buy than rent.
- At what point you should stop renting and buy?



- There is some number t of ski trips that you will take before stopping
- Suppose you are told t in advance
- **Then it is easy to decide: rent or buy**
- □ If tx<=y, then rent otherwise buy right at the start
- OFFLINE ski-rental problem
- Its solution is called the OPTIMAL SLOUTION and the cost of optimal solution is called OPTIMAL COST
- Optimal cost is tx for t<=c and y for t>c,



In the online version of the problem, the rent or buy decision must be made prior to each trip, without knowledge of t

- Strategy: rent until c=y/x trips have occurred, and then buy if a (c+1)st trip happens
- How well this strategy would do?



- If t<=c, then it is optimal minimum possible amount is spent
- If t>c, then the cost is exactly twice the optimal cost!
- The strategy can be optimal for some situations and in the worst case it incurs a cost that is twice the optimal cost
- This worst case ratio between the cost incurred by the online strategy and the optimal cost is called the 'COMPETITIVE RATIO'



- Is there a better strategy given the rules of the game?
- A strategy is simply a value 'k': the number of times to rent before buying
- **Cost of strategy:**
 - tx for t<=k
 - kx + y for t>k
- Clearly, there is no value of k that is guaranteed to achieve optimal cost in all cases



- Any k is non-optimal for the case t=k+1
- **Optimal cost = tx = (k+1)x**
- Online cost = kx +y
- $\Box kx+y = kx+cx >= (k+2)x > (k+1)x = tx$
- **This is typical of online problems**
- Without future knowledge, there is no online algorithm that is <u>always</u> optimal



- It is not hard to see that no strategy can have a competitive ratio that is less that 2
- The worst case ratio between the online cost and the optimal cost is

kx+y/[min(tx,y)] OR
max(kx+y/tx, kx+y/y)

- If k=0, then for t=1, first ratio is y/x which by assumption is at least 2
- If kx<=y, then the ratio is at least 2 when t=k (first ratio in the max)</p>
- If kx>y, then the ratio is at least 2 when t>k (second ratio in the max)



Renting costs \$20 a day
 Buying costs \$300



Omniscient strategy (if you know in advance you will ski x days:

- If x < 15, optimal policy is to rent.</p>
- If x > 15, optimal policy is to buy the <u>first day</u>.
- If x = 15, both policies are the same.

An Online strategy is described by a threshold z:

Rent for up to z days, then buy, if still skiing.



- **Offline Solution**
- If Tamon knew today that he would be skiing for d days (Instance I_d), his problem is easy
- □ If 20d <= 300 then rent
 - Else buy
- Offline optimum cost
 - OPT $(I_d) = min (20d, 300)$
 - ...BUT Tamon does not know d!!



- General Online Ski Rental Algorithm A_x
- Rent for up to x days
- Then buy, if still skiing

How to evaluate the cost of an online algorithm?



General Online Ski Rental Algorithm A_x (Rent for x days, then buy) If Tamon ends up skiing d days, his actual cost is $C(A_x, I_d) = \begin{bmatrix} 20d & if d < x \end{bmatrix}$ 20x + 300 otherwise

Whereas he could have only paid OPT (I_d) = min (20d, 300) ... but we don't know which case will apply!

GENERALIZATION OF SKI RENTAL PROBLEM



- Ski rental is relevant not only to the management of sports equipment
- Applicable to wide variety of resource allocation problems
- **Gamma** For example: power management in a laptop computer
- Laptop powers down the hard drive when it isn't in use, because running a hard drive consumes battery power
- It takes significant amount of power and time, however, to restart the hard drive
- If the user of the laptop doesn't use the hard drive for a while, how long the laptop should wait to powering it down?
- A typical online problem!!